

APPARATUS AND METHOD FOR SYNCHRONIZING PRESENCE
ATTRIBUTE DATA BETWEEN TERMINAL AND SERVER

PRIORITY

This application claims priority to an application entitled “APPARATUS
5 AND METHOD FOR SYNCHRONIZING PRESENCE ATTRIBUTE DATA
BETWEEN TERMINAL AND SERVER”, filed in the Korean Intellectual
Property Office on October 17, 2003 and assigned Serial No. 2003-72645, the
contents of which are incorporated herein by reference.

BACKGROUND OF THE INVENTION

10 **1. Field of the Invention**

The present invention relates to an apparatus and method for
synchronizing data between a terminal and a server, and more particularly to an
apparatus and method for synchronizing presence attribute data when the
terminal and the server are connected to each other to establish a messenger
15 service between them.

2. Description of the Related Art

Typically, if a client terminal is connected to a server to establish data
communication between them, there is a need for data of the client terminal to be
synchronized with that of the server. Fig. 1 shows an exemplary system for

connecting the client terminal with the server. The client terminals 100, 102, and 104 are terminals for providing a messenger service, and are connected to a server 110 over the Internet 108 or a mobile communication wireless network 106. The server 110 for providing such a messenger service includes a
5 presence attribute database (DB) 112. The presence attribute DB 112 stores various presence attributes, for example, friend list information of a corresponding client terminal 10, and status, address and contact information of individual friends contained in the friend list, etc. The presence attributes are various data stored in the client terminal 10 and the server 110 to establish a
10 messenger service. The client terminal 10 gains access to the server 110, requests presence attributes from the server 110 to receive the messenger service, and receives the presence attributes from the server 110.

Fig. 2 depicts a conventional process for synchronizing presence attribute data between the client terminal 10 and the server 110.

15 Referring to Fig. 2, upon receiving a messenger service request command from a user, the client terminal 10 transmits a login request signal to the server 110 at step 204. The server 110 having received the login request signal recognizes the client terminal 10 having transmitted the login request signal, and transmits a response signal (not shown) to the login request signal to the client
20 terminal 10, such that user service information and client terminal 10's information can be transmitted between the client terminal 10 and the server 110.

Thereafter, the client terminal 10 requests presence attributes from the server 110, and the server 110 reads presence attributes at step 206 of the client terminal 10 from the presence attribute DB 112, and transmits at step 208 the read

presence attributes to the client terminal 10.

For example, if a call connection state between the server 110 and the client terminal 10 is completed at 3 o'clock, and the client terminal 10 re-accesses the server 110 at 4 o'clock, there is no way to verify the presence attributes of the client terminal 10 during the time between 3 o'clock and 4 o'clock, such that the server 110 must verify all data of the client terminal 10. Due to this problem, the server 110 must unnecessarily verify even that data which was unchanged during the one hour. In more detail, if the client terminal 10 is disconnected from the server 110 at 3 o'clock, and is connected to the server 110 at 4 o'clock, the client terminal 10 must request all the presence attributes from the server 110, and thereby the server 110 must unavoidably transmit the same presence attributes as those at 3 o'clock to the client terminal 10 even though the client terminal 10's presence attributes stored in the server 110 are the same at 3 and 4 o'clock. In the case where a small amount of update attributes are contained in the presence attributes associated with the client terminal 10 at 3 and 4 o'clock, the server 110 must transmit all the presence attributes to the client terminal 10, such that the amount of traffic is increased and a long period of time is consumed due to such unnecessary data transmission.

As stated above, the client terminal 10 requests all of its own presence attributes from the server 110 whenever it gains access to the server 110, and receives the requested presence attributes from the server 110 in such a way that the client terminal 10 updates its own presence attributes. For example, if specific presence attributes generated when the client terminal 10 is disconnected

from the server 110 are equal to those generated when the client terminal is reconnected to the server 110, or if a small amount of updated presence attributes are provided, the server 110 must transmit all the presence attributes to the client terminal 10, such that unnecessary traffic and time are consumed.

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SUMMARY OF THE INVENTION

Therefore, the present invention has been made in view of the above problems, and it is an object of the present invention to provide an apparatus and method for controlling a server to limit the transmission of presence attribute data to only that presence attribute data that needs to be updated to a client
10 terminal when the client terminal is reconnected to the server, resulting in more effective data synchronization.

It is another object of the present invention to provide an apparatus and method for controlling the server to limit the transmission of presence attribute data to only that presence attribute data that needs to be updated to the client
15 terminal when the client terminal is reconnected to the server, thereby preventing traffic from being increased due to unnecessary data transmission.

It is yet another object of the present invention to provide an apparatus and method for controlling the client terminal to receive only that presence attribute data that needs to be updated from the server when the client terminal is
20 reconnected to the server, thereby changing current presence attribute data to effective or valid data within a short period of time.

In accordance with one aspect of the present invention, the above and other objects can be accomplished by the provision of an apparatus for providing a messenger service, comprising: a client terminal for requesting presence attribute data to be updated from a server when it is reconnected to the server
5 having been previously connected to the client terminal to perform the messenger service, and updating presence attribute data stored during a previous connection upon receiving the presence attribute data to be updated from the server; and a server for receiving a presence attribute data request to be updated from the client terminal, and transmitting only updated presence attribute data, created
10 after releasing the previous connection, from among current presence attribute data to the client terminal.

In accordance with another aspect of the present invention, there is provided a method for synchronizing presence attribute data between a client terminal and a server in an apparatus composed of the client terminal and the
15 server providing the client terminal with a messenger service, comprising the steps of: if the client terminal is reconnected to the server having been previously connected to the client terminal to perform the messenger service, controlling the client terminal to request presence attribute data to be updated from the server; upon receiving a presence attribute data request to be updated from the client
20 terminal, controlling the server to transmit only updated presence attribute data, created after releasing the previous connection, from among current presence attribute data to the client terminal; and controlling the client terminal to update presence attribute data stored at a time of a previous connection upon receiving presence attribute data to be updated from the server.

BRIEF DESCRIPTION OF THE DRAWINGS

The above and other objects, features and other advantages of the present invention will be more clearly understood from the following detailed description taken in conjunction with the accompanying drawings, in which:

5 Fig. 1 is a block diagram illustrating a conventional system for establishing communication between a client terminal and a server;

 Fig. 2 is a flow chart illustrating a conventional process for synchronizing presence attribute data between the client terminal and the server shown in Fig. 1;

10 Fig. 3 is a flow chart illustrating a process for synchronizing presence attribute data between a client terminal and a server in accordance with a preferred embodiment of the present invention;

 Fig. 4 is a flow chart illustrating a process for controlling the client terminal to synchronize its presence attribute data with the server in accordance
15 with a preferred embodiment of the present invention; and

 Fig. 5 is a flow chart illustrating a process for controlling the server to synchronize its presence attribute data with the client terminal in accordance with a preferred embodiment of the present invention.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS

20 Now, preferred embodiments of the present invention will be described in detail with reference to the annexed drawings. In the drawings, the same or similar elements are denoted by the same reference numerals even though they

are depicted in different drawings. In the following description, a detailed description of known functions and configurations incorporated herein will be omitted when it may obscure the subject matter of the present invention.

At the moment of disconnecting a client terminal 10 from a server 110,
5 the server 110 and the client terminal 10 store the same data therein. If data is created and stored when the server 110 is disconnected from the client terminal 10, and then the server 110 is reconnected to the client terminal 10 using a specific key value for identifying the data, the client terminal 10 transmits only the specific key value to the server 110, and the server 110 receives the specific
10 key value from the client terminal 10. In this case, provided that the same key value as the specific key value and data corresponding to the key value are stored in the server, the server 110 transmits the data that needs to be updated to the client terminal 10.

An apparatus for establishing communication between the client terminal
15 10 and the server 110 according to an embodiment of the present invention will hereinafter be described with reference to Fig. 1.

The client terminal 10 includes all the terminals capable of providing a messenger service, for example, a computer 100, a mobile phone or terminal 102, and a PDA (Personal Digital Assistant) 104, etc. If the client terminal 10 is
20 disconnected from the server 110, it stores presence attribute data, such as a friend list data and friend condition data, etc., created during the last communication between the client terminal 10 and the server 110. The client terminal 10 stores a unique session ID of a previous session between the client

terminal 10 and the server 110, a client ID of the client terminal 10 storing the friend list and condition data, and a transaction ID between the client terminal 10 and the server 110. In this case, the transaction designates a response of a server 110 to a request of a client terminal 10, or a response from a client
5 terminal 10 to a request by a server. Whenever a session between the server 110 and the client terminal 10 is established, a unique transaction ID is assigned to either the client terminal 10 or the server 110. In this case, during a communication time between the server 110 and the client terminal 10, the client terminal 10 stores a normally-completed transaction ID generated when request
10 and response operations between the server 110 and the client terminal 10 are all normally completed.

In the case of reconnecting the client terminal 10 to the server 110, the client terminal 10 creates a key composed of such a session ID, a client ID of the client terminal 10, and a transaction ID corresponding to the last communication
15 with the server 110, and transmits the created key to the server 110.

The server 110 connected to the client terminal 10 over the Internet provides the client terminal 10 with a messenger service when it receives a messenger service request from the client terminal 10.

The server 110 stores data for providing individual client terminals 10
20 with a messenger service in the presence attribute DB 112. The server 110 transmits presence attribute data, for example, an ID and password of a user, a client ID of the client terminal 10, a friend list defined by a session ID indicative of a connection state with the client terminal 10, and a friend condition, etc.,

created during the last communication associated with the client terminal 10 in the presence attribute DB 112. The server 110 stores various data, for example, a session ID equal to the last session ID created before it is connected to the client terminal 10, a client ID of the client terminal 10 storing the friend list and
5 condition information, and the last transaction ID of a predetermined transaction normally completed while it communicates with the client terminal 10, in the presence attribute DB 112. Upon receiving a re-access request from the client terminal 10, the server 110 determines whether key values stored in the presence attribute DB 112 are equal to the session ID stored in the received re-access
10 request, the client ID of the client terminal 10, and the transaction ID corresponding to the last communication with the server 110.

If the presence attribute DB 112 stores the same key value as a key value created when the client terminal 10 is reconnected to the server 110, the server 110 compares presence attribute data with the last valid presence attribute data
15 stored in the presence attribute DB 112 according to the same key value at a time of a connection release between the server 110 and the client terminal 10, and reads a different field between both presence attribute data. The server 110 inserts the different field in a response data field of the presence attribute request, and transmits it to the client terminal 10. In this case, if there is no update
20 information between current presence attribute data and previous presence attribute data stored in response to the same key value, the server 110 transmits the response data field with the null status. After presence attribute data of the server 110 is synchronized with that of the client terminal 10, the server 110 may delete the key value used for such synchronization and presence attribute data
25 corresponding to the key value.

If the presence attribute DB 112 does not store the same key value as a key value created at a re-access time between the server 110 and the client terminal 10, the server 110 transmits a specific message indicative of a presence attribute synchronization failure to the client terminal 10. The client terminal
5 10 receiving the message deletes its own stored key and associated data because they are no longer effective. The client terminal 10 requests presence attributes from the server 110 using a general presence synchronization process, receives all the current presence attributes from the server 110, updates the received presence attributes, and synchronizes them with the server 110.

10 The above method for synchronizing presence attribute data at a re-access time between the client terminal 10 and the server 110 will hereinafter be described with reference to Fig. 3.

If the client terminal 10 is connected to the server 110, the client terminal 10 and the server 110 share a session ID and a client ID associated with such
15 connection. In this case, if the client terminal 10 is disconnected from the server 110, the client terminal 10 and the server 110 each store a successful transaction ID created during the last communication between them. If the client terminal 10 is normally disconnected from the server 110, the transaction ID stored in the client terminal 10 and the server 110 adapts a logout request or a
20 connection end request as the last successful transaction. Otherwise, if the client terminal 10 is abnormally disconnected from the server 110, a transaction just before the moment of a connection failure is adapted as the last successful transaction. According to the present invention, it is assumed that the client

terminal 10 and the server 110 each share presence attribute data, and are then disconnected from each other. If the client terminal 10 and the server 110 are disconnected from each other, the client terminal 10 and the server 110 each create a synchronization key for synchronizing presence attribute data composed
5 of a session ID, a client ID, and a transaction ID, and store the created synchronization key. The client terminal 10 and the server 110 store such a synchronization key and presence attribute data .

If the client terminal 10 transmits a re-access or login request to the server 110 at step 300 after being disconnected from the server 110, the server 110
10 recognizes the client terminal 10 transmitting the re-access request, and transmits a response signal to the re-access request to the client terminal 10, such that user service information and the client terminal 10 information are communicated between the client terminal 10 and the server 110. Thereafter, at step 304, the client terminal 10 transmits a predetermined message determining whether the
15 server 110 provides a specific service for synchronizing presence attributes using the above synchronization key (i.e., synckey) to the server 110. If it is determined that the server 110 provides the specific service for synchronizing the presence attributes using the synchronization key, the server 110 transmits a predetermined response message to the client terminal 10 at step 306.

20 If the client terminal 10 receives a response message to a presence attribute synchronization service request at step 306, it requests at step 308 the presence attribute data from the server 110. In this case, the client terminal 10 transmits the synchronization key (synckey) to synchronize presence attribute data composed of a session ID, a client ID, and a transaction ID in such a way

that it requests presence attributes from the server 110.

The server 110 determines at step 310 whether the same synchronization key (synckey) as the synchronization key (synckey) received from the client terminal 10 is stored in the presence attribute DB 12. If it is determined that the same synchronization key (synckey) is stored in the presence attribute DB 112, the server 110 goes to step 314. Otherwise, if the same synchronization key (synckey) is not stored in the presence attribute DB 112, the server 110 transmits at step 312 a specific message indicative of a presence attribute synchronization failure to the client terminal 10. The client terminal 10 receiving the message deletes its own stored key and associated data because they are no longer effective. The client terminal 10 requests presence attributes from the server 110 using a general presence synchronization process, receives all the current presence attributes of friends contained in the user's friend list from the server 110, updates the received presence attributes, and synchronizes them with the server 110.

If it is determined in step 310 that the same key value as the key value received from the client terminal 10 is stored in the presence attribute DB 112, the server 110 at step 314 compares presence attribute data indicative of the current condition of the current presence attribute data with the last valid presence attribute data stored in the presence attribute DB 112 according to the same key value at a time of a connection release between the server 110 and the client terminal 10, and reads an updated field between both presence attribute data. The server 110 inserts the updated field in a response data field of the presence attribute request, and transmits it to the client terminal 10 at step 316.

In this case, if there is no update information between current presence attribute data and previous presence attribute data stored in response to the same key value, the server 110 transmits the response data field with the null status. After presence attribute data of the server 110 is synchronized with that of the client terminal 10, the server 110 may delete the key value used for such synchronization and presence attribute data corresponding to the key value.

As shown in Fig. 3, in the case of reconnecting the client terminal 10 to the server 110 using the synchronization key value to synchronize presence attributes of the client terminal with those of the server 110, the server 110 may select only current updated presence attributes from among all presence attributes stored at a time of a connection release, and may transmit the selected presence attributes to the client terminal 10.

A method for controlling the client terminal 10 to synchronize its presence attribute data with that of the server 110 will hereinafter be described with reference to Fig. 4. Fig. 4 is a flow chart illustrating a process for controlling the client terminal 10 to synchronize its presence attribute data with the server 110 in accordance with a preferred embodiment of the present invention. **[PLEASE CORRECT FIG. 4 AS SHOWN.]**

Referring to Fig. 4, if a process goes from step 400 indicative of a standby mode to step 402 and a user enters a login request to the server 110 to receive a messenger service at step 402, the client terminal 10 transmits the login request to the server 110 at step 404. If the client terminal 10 receives a response message to the login request from the server 110 at step 406, user service

information and the client terminal 10's information are transmitted to the server 110 at step 408. The client terminal 10 transmits a prescribed request to the server 110 such that the server 110 determines whether a presence attribute synchronization service function can be executed or not using the synchronization key (synckey) value at step 410. If the client terminal 10 receives a response message indicating that the presence attribute synchronization service function can be executed using the synchronization key (synckey) value at step 412 from the server at step 412, it goes to step 414. Otherwise, if the client terminal 10 does not receive the response message at step 412, it performs a general presence synchronization process at step 422. During the general presence synchronization process, the client terminal 10 requests presence attributes from the server 110, receives all the current presence attributes from the server 110, updates the received presence attributes, and synchronizes them with the server 110.

After performing the above step 412, the client terminal 10 creates a synchronization key (synckey) composed of a session ID, a client ID, and a transaction ID at step 414, includes the created synchronization key (synckey) in a presence attribute synchronization request message, and transmits the presence attribute synchronization request message to the server 110 at step 416. Upon receiving presence attribute data to be updated from the server 110 at step 418, the client terminal 10 at step 420 updates only the received presence attributes from among the stored presence attributes, and stores the updated presence attributes. If no updated presence attribute data is received at step 418, the client terminal 10 performs a general presence attribute synchronization at step 422.

If the client terminal 10 requests presence attributes from the server 110 using the synchronization key value, the server 110 transmits presence attributes to the client terminal 10 such that presence attributes of the server 110 are synchronized with those of the client terminal 10, its detailed description will
5 hereinafter be described with reference to Fig. 5. Fig. 5 is a flow chart illustrating a process for controlling the server 110 to synchronize its presence attribute data with the client terminal 10 in accordance with a preferred embodiment of the present invention.

Referring to Fig. 5, the server 110 determines at step 502 whether a login
10 request is received from the client terminal 10 while maintaining a standby mode at step 500. Upon receiving the login request from the client terminal 10, the server 110 transmits a response message to the login request at step 504. If the server 110 receives a prescribed request from the client terminal 10 to determine whether a presence attribute synchronization service can be executed using a
15 synchronization key (synckey) value at step 506, the process goes to step 508. Otherwise, if the server 110 does not receive the prescribed request from the client terminal 10 at step 506, the process goes to step 522. If the server 110 goes to step 508, it determines whether a presence attribute synchronization service can be executed using the synchronization key (synckey) value at step
20 508. If it is determined that the presence attribute synchronization service can be executed at step 508, the server goes to step 510. Otherwise, if it is determined that the presence attribute synchronization cannot be executed at step 508, the server 110 goes to step 520, such that at step 520 it transmits a prescribed message indicating that it is impossible to perform the presence

attribute synchronization service using the synchronization key value to the client terminal 10, and terminates overall processes. If it is determined that the presence attribute synchronization service can be executed using the synchronization key value at step 508, the server 110 transmits at step 510 a
5 response message to the request for checking the presence attribute synchronization service enable status to the client terminal 10. If the client terminal 10 requests presence attributes containing the synchronization key value from the server 110 at step 512, the server 110 determines whether a synchronization key value equal to a reception synchronization key value is
10 stored in the presence attribute DB 112 at step 514. If the values are not equal, a synchronization failure message is sent in step 526. If it is determined that the same synchronization key value as the reception synchronization key value is found at step 514, the server 110 at step 516 compares current presence attributes with the last presence attributes in response to the synchronization key value, and
15 reads updated presence attributes. The server 110 transmits presence attribute data to be updated to the client terminal 10 at step 518. The server 110 compares presence attribute data with the last valid presence attribute data stored in response to the same key value at a time of a connection release between the server 110 and the client terminal 10, and reads an updated field
20 between both presence attribute data. For example, provided that a connection release between the server 110 and the client terminal 10 occurs on the condition that three friends are in an online status and the remaining friends other than the three friends are in an offline status, the client terminal 10 and the server 110 store presence attribute data at a time of the connection release, and store a
25 synchronization key value composed of a client ID, a session ID, and a transaction ID. The client terminal 10 inserts a synchronization key value in the

presence attributes when it is reconnected to the server 110, and requests presence attributes from the server 110. If a connection release between the server 110 and the client terminal 10 occurs, the server 110 receiving the presence attributes compares current friend status with a specific status in which
5 three friends contained in the friend list serving as stored presence attributes is in the online mode and the remaining friends other than the three friends are in the offline mode. In more detail, if a connection release between the server 110 and the client terminal 10 occurs, the server 110 determines whether an online friend has changed to an offline friend, and determines whether status
10 information (e.g., an ID and an address, etc.) has changed to other information in such a way that it transmits only changed information to the client terminal 10. However, if such a connection release between the server 110 and the client terminal 10 occurs, and there is no update information between the stored presence attributes and new presence attributes created at a re-connection time
15 between the server 110 and the client terminal 10, the server 110 transmits data indicative of a null status to a response data field such that it informs the client terminal 10 of no updated presence attributes. If a synchronization process is completed, the server 110 may delete any used synchronization key value and its associated presence attribute data.

20 After performing steps 506 and 512 and a “no” is determined, the server 110 determines whether a general presence attribute synchronization request is received from the client terminal 10 at step 522. If it is determined that the client terminal 10 transmits the general presence attribute synchronization request, the server 110 transmits at step 524 all the current presence attribute
25 data to the client terminal 10, and terminates overall operations.

As described above, if the client terminal 10 is reconnected to the server 10 to perform a messenger service, the inventive apparatus receives not all the presence attributes but only updated presence attributes from among presence attributes stored at a time of a connection release, such that it can quickly
5 synchronize the client terminal 10 with the server 110 at a connection moment between the client terminal 10 and the server 110, and can reduce the amount of traffic due to transmission/reception of the small amount of data.

As apparent from the above description, if the client terminal is reconnected to the server, there is no need for the client terminal and the server
10 to mutually exchange specific data that need not be synchronized with the server, resulting in reduction of the amount of transmission data between the client terminal and the server. Furthermore, if the client terminal is connected to the server, current status data of friends contained in the user's friend list can be changed to effective or valid data within a short period of time, resulting in a
15 reduced synchronization time.

Although the preferred embodiments of the present invention have been disclosed for illustrative purposes, those skilled in the art will appreciate that various modifications, additions and substitutions are possible, without departing from the scope and spirit of the invention as disclosed in the accompanying
20 claims.